

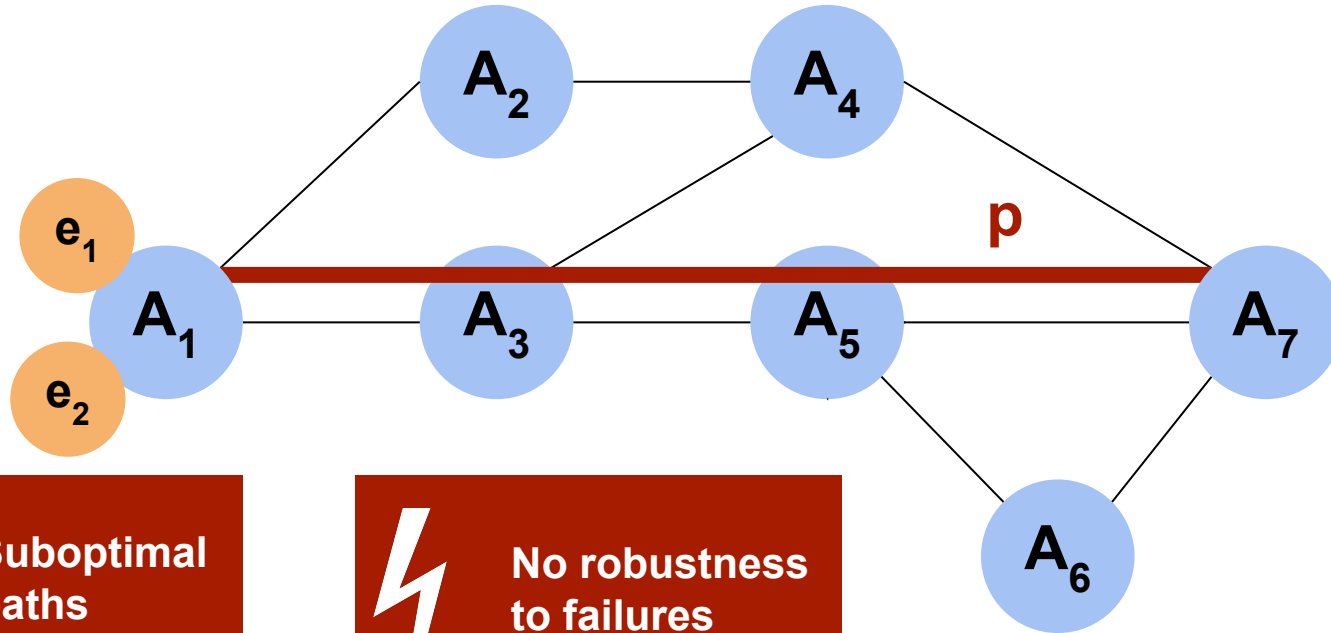


The Value of Information in Selfish Routing

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Network-based path selection



Suboptimal
paths

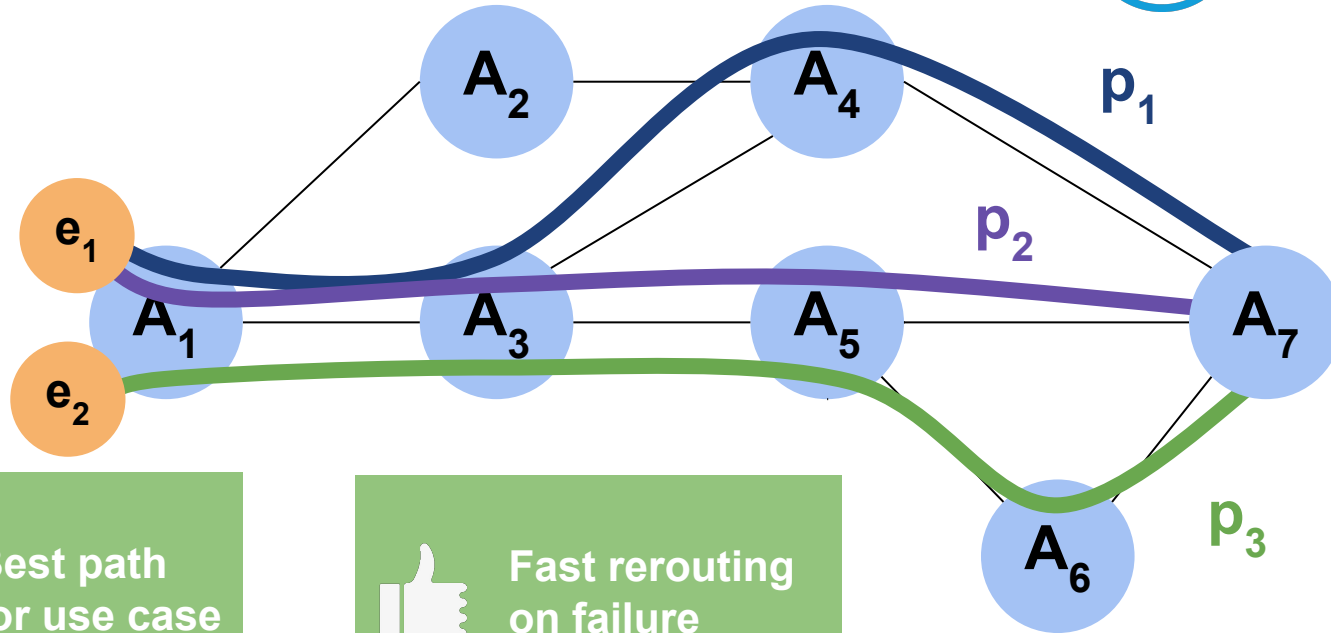


No robustness
to failures

Source-based path selection



SCION

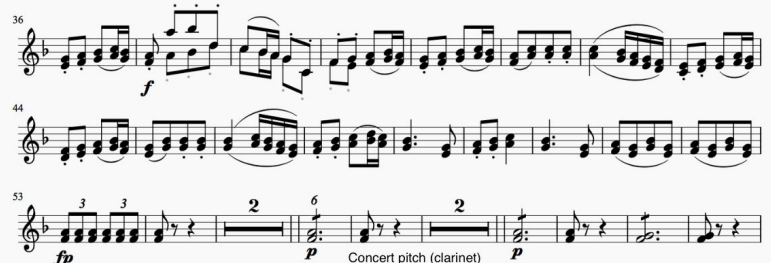


Best path
for use case

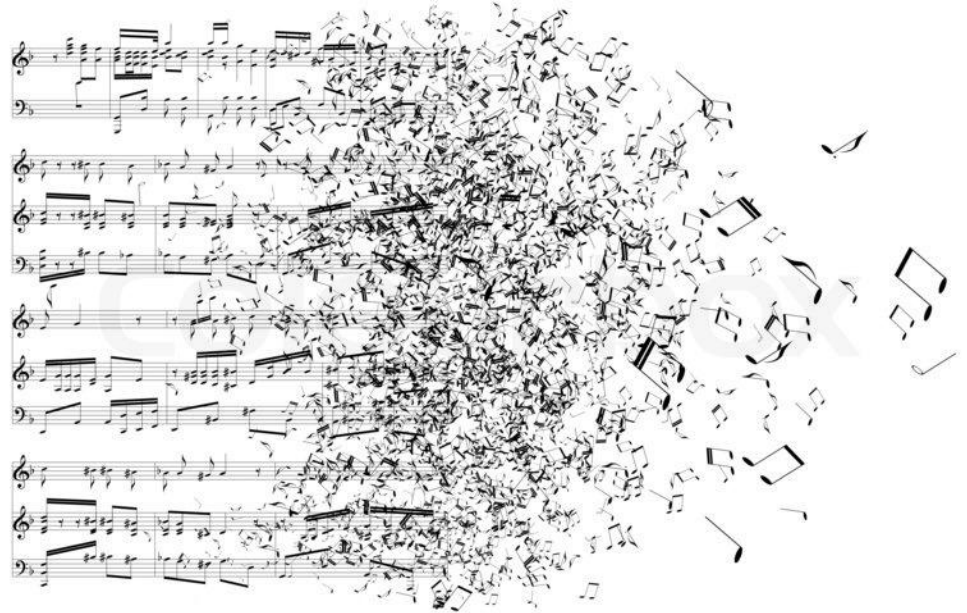


Fast rerouting
on failure

Network-based path selection: Network operator view



Source-based path selection: Network operator view



Goals of our work

Revisit selfish-routing concepts to investigate two issues arising in emerging path-aware Internet architectures:

- **Impact of information:** What network state information should be shared with end-hosts?
- **Impact on network operators:** What is the impact of selfish routing on the cost of network operators?

Price of Anarchy: Three components

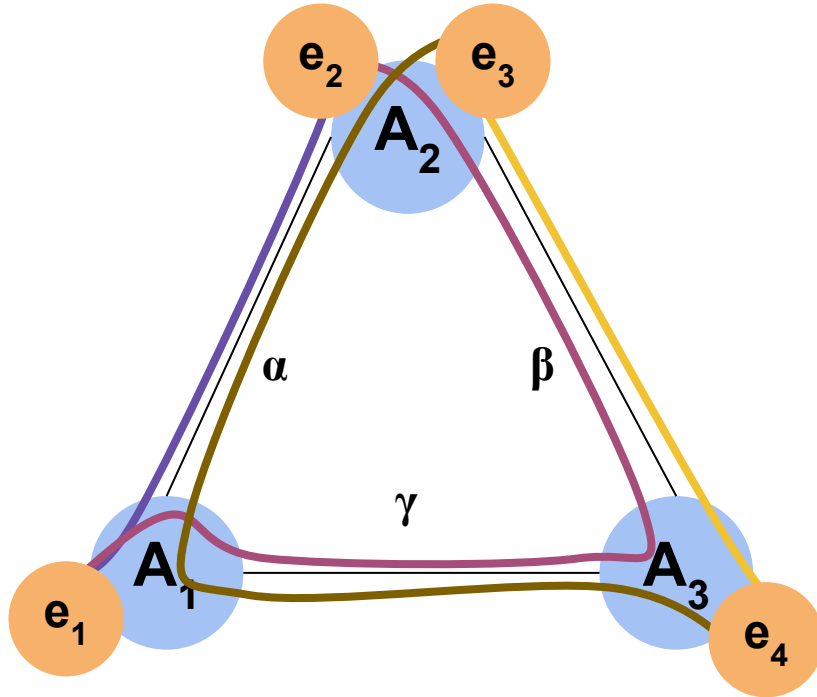
C Social cost function

F^{opt} Social optimum

F^{eq} Equilibrium

$$\text{PoA} = \frac{C(\mathbf{F}^{\text{eq}})}{C(\mathbf{F}^{\text{opt}})}$$

Adapted Wardrop model of source-based path selection



$$\mathbf{d} = (d_{1,2}, d_{3,4}) = (1, 1)$$

$$\mathbf{F} = (F_{\alpha}, F_{\gamma\beta}, F_{\beta}, F_{\alpha\gamma})$$

$$\mathbf{f} = (f_{\alpha}, f_{\beta}, f_{\gamma})$$

$$c_{\alpha}(f_{\alpha}) = 1$$

$$c_{\beta}(f_{\beta}) = f_{\beta}^2$$

$$c_{\gamma}(f_{\gamma}) = f_{\gamma}$$

$$C_{\pi}(\mathbf{F}) = \sum_{\ell \in \pi} c_{\ell}$$

Total cost functions and social optima

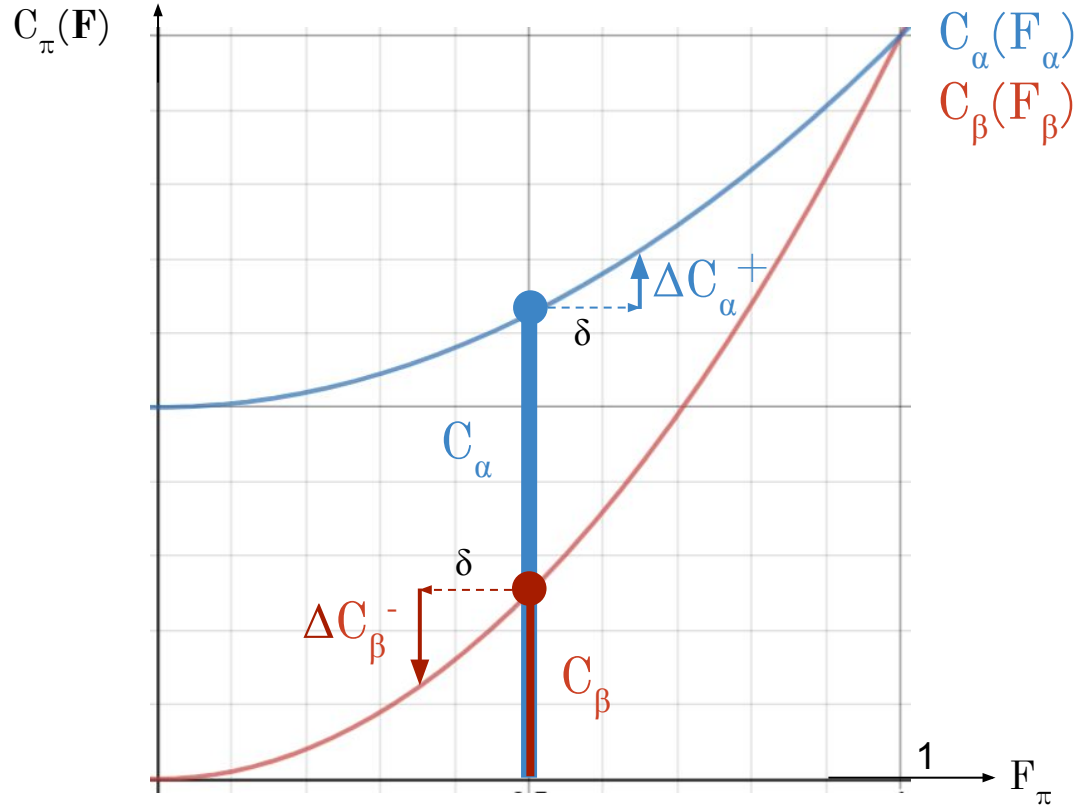
End-host cost function: $C^* = \sum_{\text{end-hosts}} \sum_{\text{paths}} \text{flow on path} \cdot \text{path cost}$
(classic) $= \sum_{\pi \in \Pi} F_{\pi} \cdot C_{\pi}(\mathbf{F}) = \sum_{\ell \in L} f_{\ell} \cdot c_{\ell}(f_{\ell})$

End-host optimum: $\mathbf{F}^* = \operatorname{argmin}_{\mathbf{F}} C^*(\mathbf{F})$

Network-operator cost function: $C^{\#} = \sum_{\text{links}} \text{link cost} = \sum_{\ell} c_{\ell}(f_{\ell})$

Network-operator optimum: $\mathbf{F}^{\#} = \operatorname{argmin}_{\mathbf{F}} C^{\#}(\mathbf{F})$

Characterizing social optima: Suboptimal path flow pattern



$$\mathbf{d} = (d_{1,2}) = (1)$$

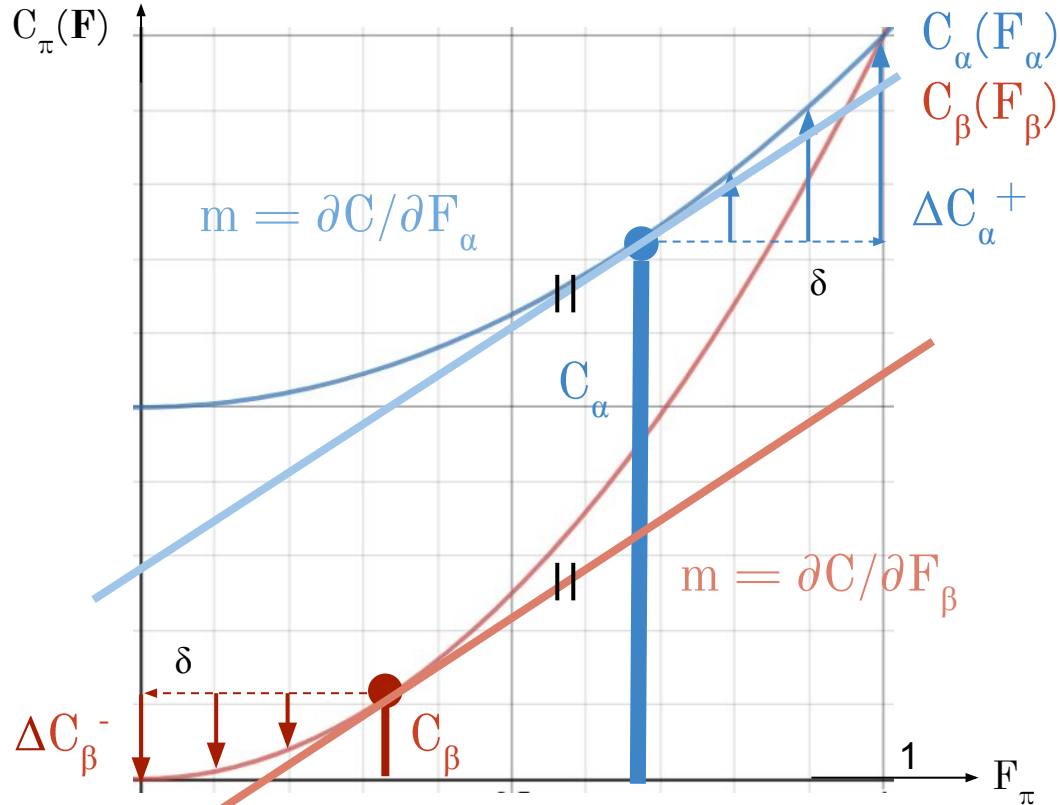
$$\mathbf{F} = (F_\alpha, F_\beta)$$

$$C(\mathbf{F}) = C_\alpha(F_\alpha) + C_\beta(F_\beta)$$

$$\exists \delta. |\Delta C_\alpha^+| < |\Delta C_\beta^-|$$

$\Rightarrow C$ can be reduced

Characterizing social optima: Optimal path flow pattern



$$\forall \delta. |\Delta C_{\alpha}^{+}| > |\Delta C_{\beta}^{-}|$$

$\Rightarrow C$ **cannot** be reduced

$\Rightarrow C$ is optimal

$$\partial C / \partial F_{\alpha} = \partial C / \partial F_{\beta}$$

$$\Rightarrow \forall \delta. |\Delta C_{\alpha}^{+}| > |\Delta C_{\beta}^{-}|$$

$\Rightarrow C$ is optimal

Socially optimal marginal costs

$\partial C(\mathbf{F})/\partial F_\pi$ is the *marginal cost* of path π

given path-flow pattern \mathbf{F}

A path-flow pattern \mathbf{F} is optimal w.r.t. a cost function $C \in \{C^*, C^\#\}$

if for every origin-destination pair:

$$F_\alpha, \dots, F_\rho > 0 \qquad F_\sigma, \dots, F_\omega = 0$$

$$\frac{\partial C(\mathbf{F})}{\partial F_\alpha} = \dots = \frac{\partial C(\mathbf{F})}{\partial F_\rho} \leq \frac{\partial C(\mathbf{F})}{\partial F_\sigma} \leq \dots \leq \frac{\partial C(\mathbf{F})}{\partial F_\omega}$$

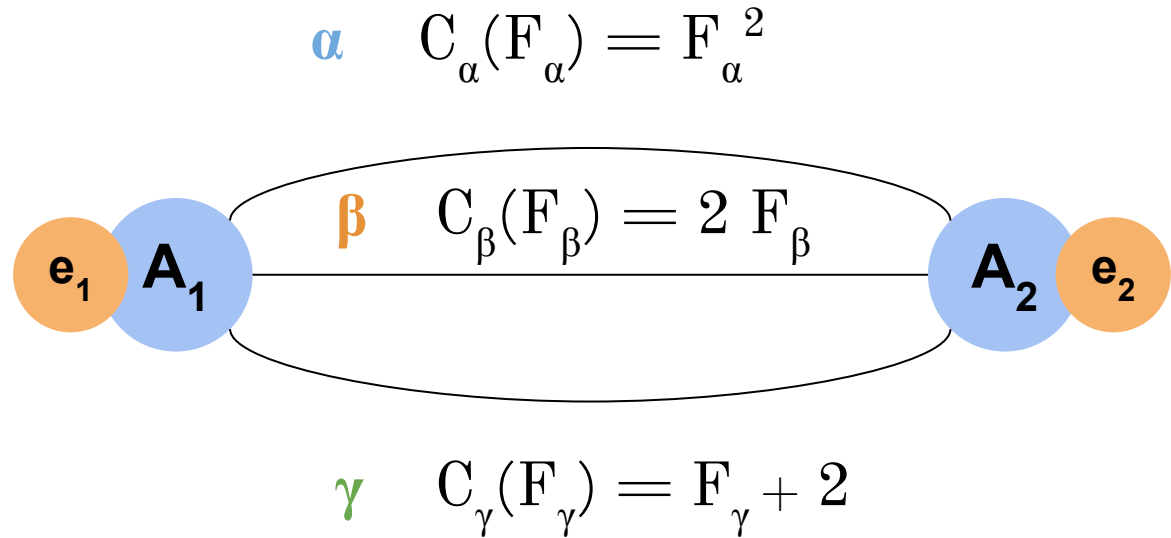
Social optimum: Comparison (Example)

$$\mathbf{F}^\# = (F_\alpha, F_\beta, F_\gamma)$$

$$= (1/2, 0, 1/2)$$

$$\mathbf{F}^* = (F_\alpha, F_\beta, F_\gamma)$$

$$= (2/3, 1/3, 0)$$



Different optima!

Network operators prefer usage of links with little variable cost (here: γ)

Price of Anarchy: Where are we?

C Total cost function

F^{opt} Social optimum

F^{eq} Equilibrium

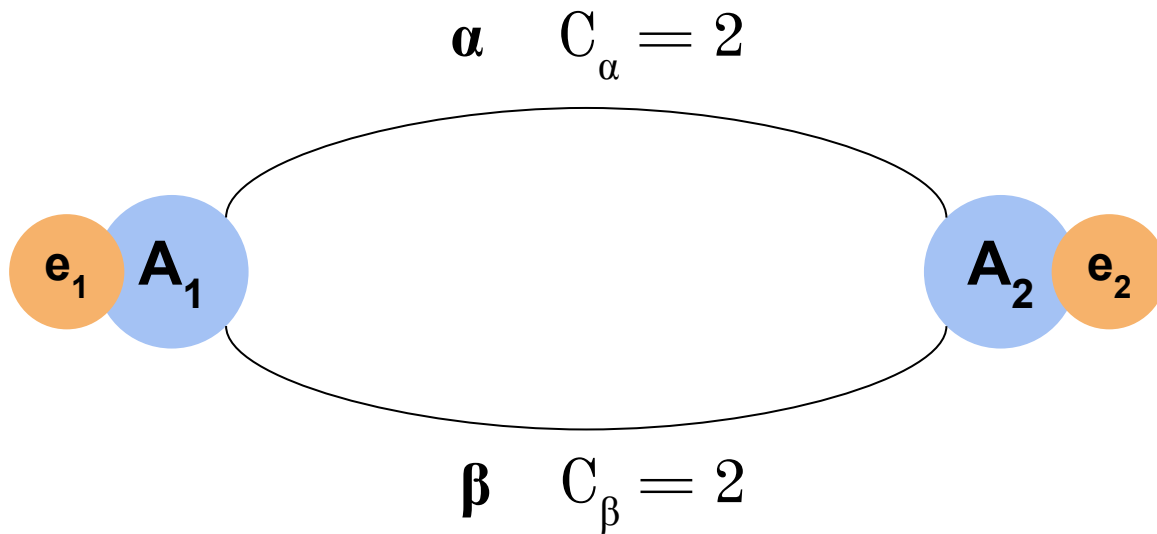


$$\text{PoA} = \frac{C(\mathbf{F}^{\text{eq}})}{C(\mathbf{F}^{\text{opt}})}$$

Equilibrium with latency-only information (LI equilibrium)

$$\mathbf{d} = (d_{1,2}) = (1)$$

$$\begin{aligned} \mathbf{F} &= (F_\alpha, F_\beta) \\ &= (1, 0) \end{aligned}$$



$$C_\alpha = C_\beta \quad \Rightarrow \quad \mathbf{F} = (1, 0) \text{ is an LI equilibrium}$$

Characterizing the LI equilibrium

A path flow pattern \mathbf{F} is an LI equilibrium
if for every origin-destination pair:

$$F_\alpha, \dots, F_\rho > 0 \qquad F_\sigma, \dots, F_\omega = 0$$

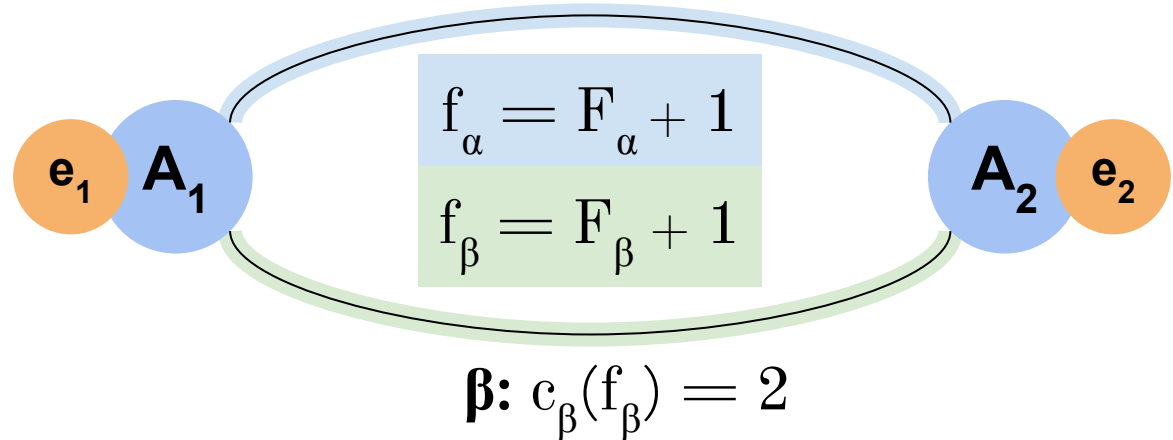
$$C_\alpha(\mathbf{F}) = \dots = C_\rho(\mathbf{F}) \leq C_\sigma(\mathbf{F}) \leq \dots \leq C_\omega(\mathbf{F})$$

Equilibrium with perfect information (PI equilibrium)

$$\alpha: c_\alpha(f_\alpha) = f_\alpha$$

$$d_{(1)} = (d_{1,2}) = (1)$$

$$F_{(1)} = (F_\alpha, F_\beta)$$



Minimize selfish cost $C_{(1)}(F_{(1)}) = F_\alpha \cdot (F_\alpha + 1) + F_\beta \cdot 2$

$$\Rightarrow (F_\alpha, F_\beta) = (2/3, 1/3) \text{ is a PI equilibrium}$$

Characterizing the PI equilibrium

A path flow pattern \mathbf{F} is a PI equilibrium

if for every origin-destination pair of any end-host e :

$$F_\alpha, \dots, F_\rho > 0 \qquad F_\sigma, \dots, F_\omega = 0$$

$$\frac{\partial C_{(e)}(\mathbf{F})}{\partial F_\alpha} = \dots = \frac{\partial C_{(e)}(\mathbf{F})}{\partial F_\rho} \leq \frac{\partial C_{(e)}(\mathbf{F})}{\partial F_\sigma} \leq \dots \leq \frac{\partial C_{(e)}(\mathbf{F})}{\partial F_\omega}$$

Capturing the value of information

**Information
assumption**

Latency-only
Information (LI)

Perfect
Information (PI)

Equilibrium

F^0

F^+

Price of Anarchy

$$PoA^0 = \frac{C(F^0)}{C(F^{opt})}$$

$$PoA^+ = \frac{C(F^+)}{C(F^{opt})}$$



Δ = Value of Information (VoI)

The benefits of information

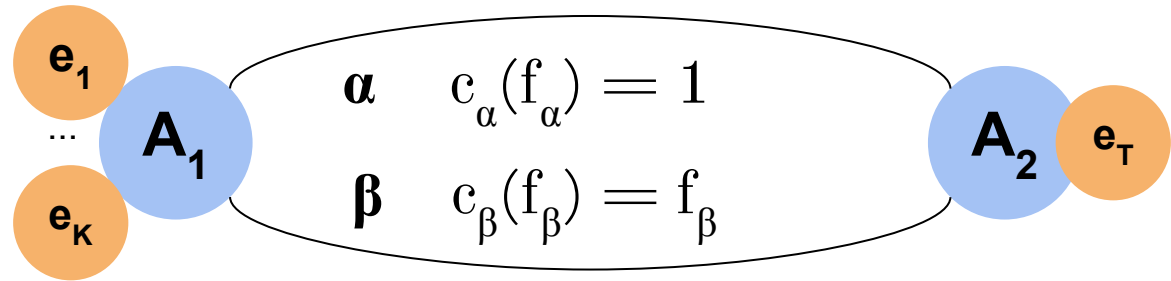
$$\text{VoI} > 0$$

The benefits of information: Network of parallel links

(cf. Roughgarden 2003)

$$\sum_k d_{k,T} = 1$$

$$\mathbf{F} = (F_{1\alpha}, F_{1\beta}, \dots, F_{K\alpha}, F_{K\beta}, \dots)$$



EH Opt: \mathbf{F}^* s.t. $f_\beta = 1/2$

LI Eq: \mathbf{F}^0 s.t. $f_\beta = 1$

NO Opt: $\mathbf{F}^\#$ s.t. $f_\beta = 0$

PI Eq: \mathbf{F}^+ s.t. $f_\beta = K/(K+1)$

The benefits of information: Network of parallel links

	LI equilibrium	PI equilibrium
End-host perspective	$\mathbf{PoA}^{*0} = \frac{4}{3}$	$\mathbf{PoA}^{*+} = \frac{(K^2 + K + 1)}{(K^2 + 2K + 1)} \cdot \frac{4}{3} \leq \mathbf{PoA}^{*0}$
Network-operator perspective	$\mathbf{PoA}^{\#0} = 2$	$\mathbf{PoA}^{\#+} = 1 + \frac{K}{(K + 1)} \leq 2 = \mathbf{PoA}^{\#0}$

The benefits of information: Network of parallel links

	LI equilibrium	PI equilibrium
End-host perspective	<p>PI equilibrium cheaper than LI equilibrium</p> <p>$V_{oI} > 0$</p>	$P_{oA}^{*+} = \frac{(K^2 + K + 1)}{(K^2 + 2K + 1)} \cdot \frac{4}{3} \leq P_{oA}^{*0}$
Network operator perspective		$P_{oA}^{\#+} = 1 + \frac{K}{(K + 1)} \leq 2 = P_{oA}^{\#0}$

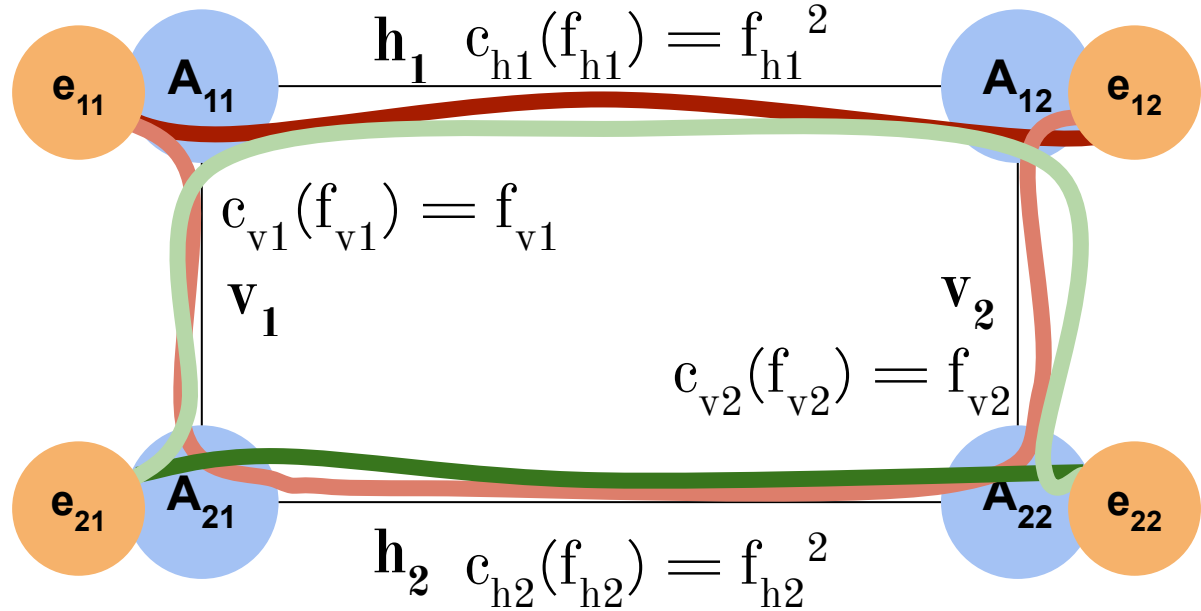
The drawbacks of information

$$\text{VoI} < 0$$

The drawbacks of information: Ladder network

$$\mathbf{d} = (d_{11,12}, d_{21,22}) \\ = (1, 1)$$

$$\mathbf{F}^{\rightarrow} = (1, 0, 1, 0)$$

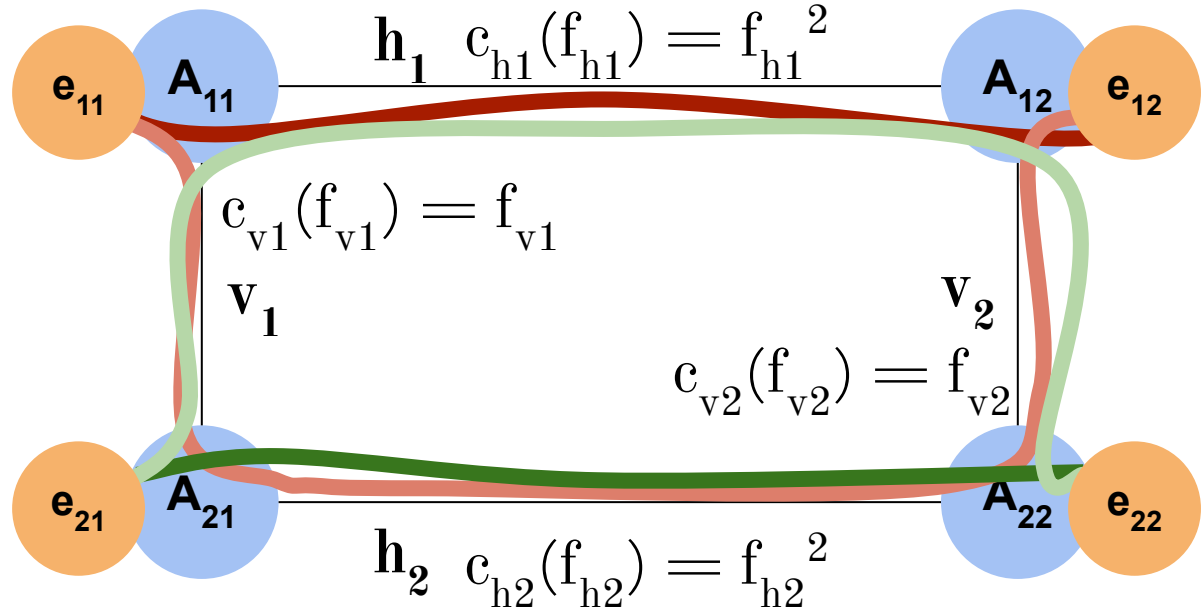


Direct-only \mathbf{F}^{\rightarrow} is universally optimal:

$$\mathbf{F}^{\rightarrow} = \mathbf{F}^* = \mathbf{F}^{\#}$$

The drawbacks of information: Ladder network

$$\mathbf{F}^{\rightarrow} = (1, 0, 1, 0)$$



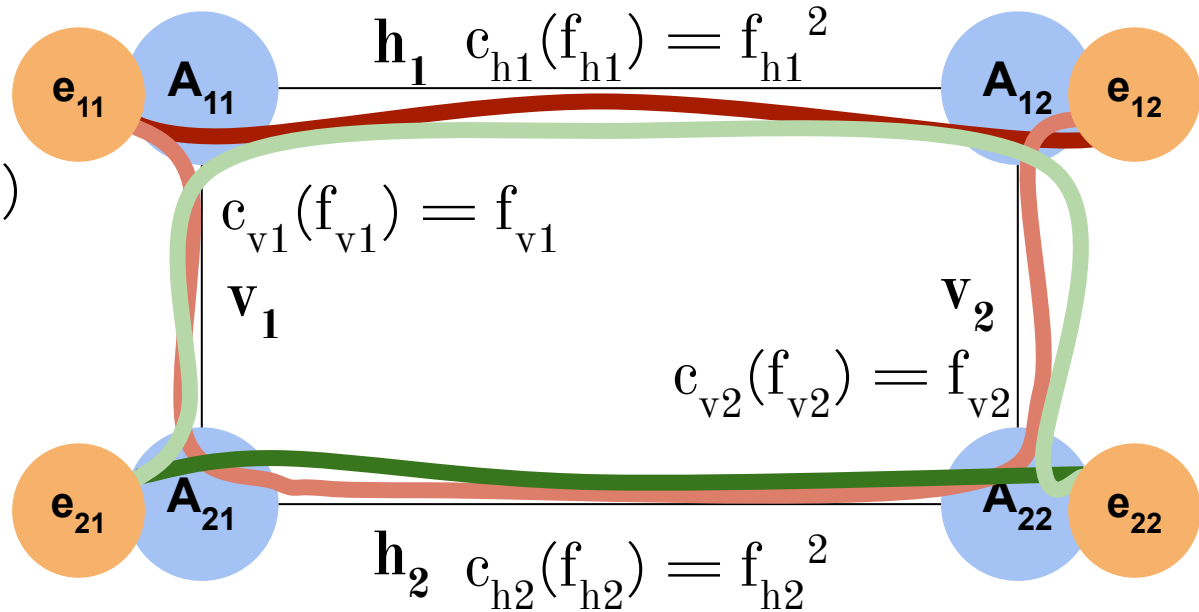
$$C_{1H}(\mathbf{F}^{\rightarrow}) = 1 = C_{1V}(\mathbf{F}^{\rightarrow}) \quad \Rightarrow \quad \mathbf{F}^{\rightarrow} = \mathbf{F}^0$$

(LI equilibrium is optimal)

The drawbacks of information: Ladder network

$$\mathbf{F}^{\rightarrow} = (1, 0, 1, 0)$$

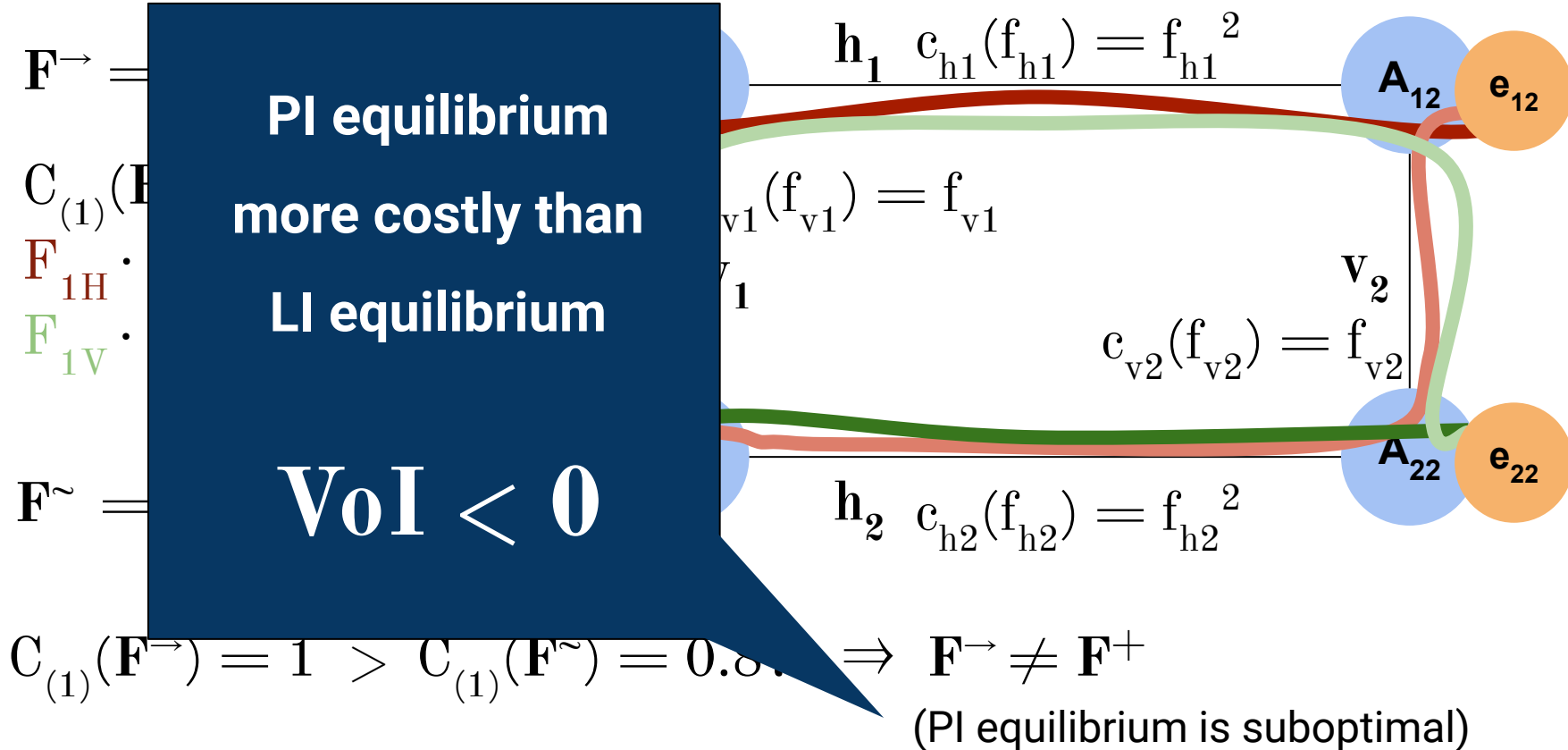
$$\mathbf{F}^{\sim} = (0.9, 0.1, 1, 0)$$



$$C_{(1)}(\mathbf{F}^{\rightarrow}) = 1 > C_{(1)}(\mathbf{F}^{\sim}) = 0.87 \Rightarrow \mathbf{F}^{\rightarrow} \neq \mathbf{F}^+$$

(PI equilibrium is suboptimal)

The drawbacks of information: Ladder network

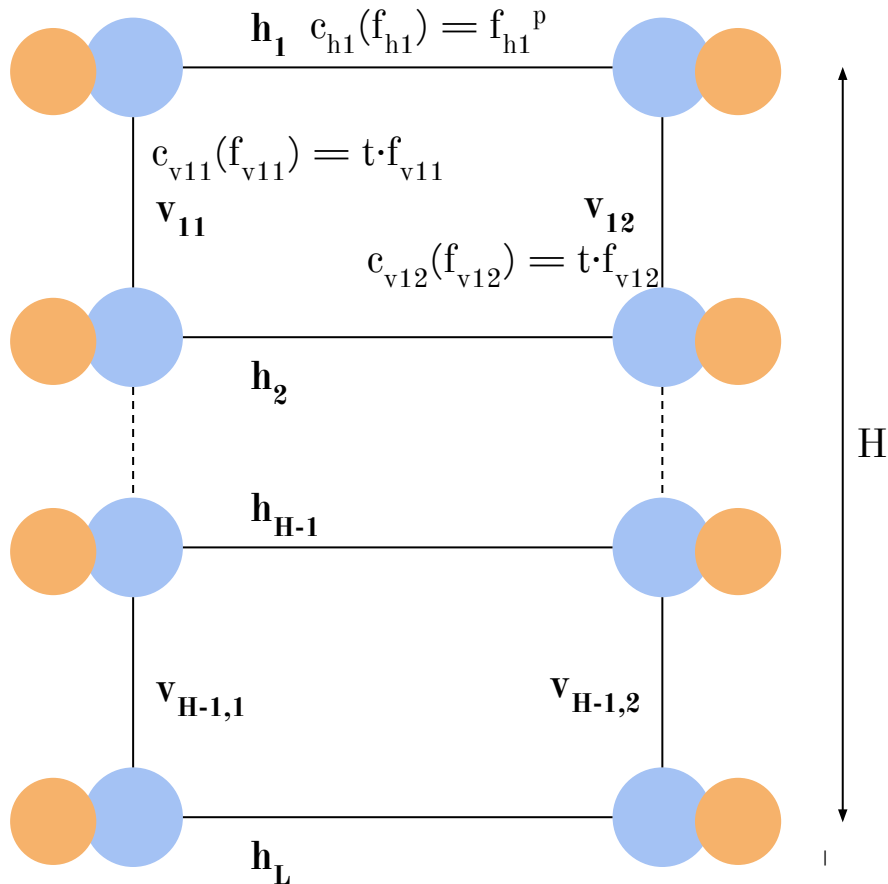


The drawbacks of information: Generalized ladder network

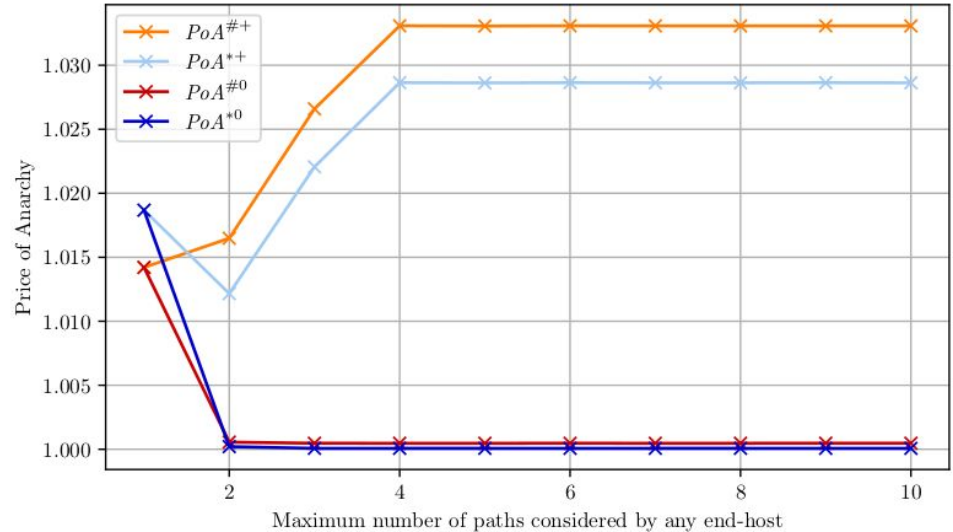
Upper bound on
PoA for network
operators:

$$\text{PoA}^{\#+} \leq 1 + \frac{2(H-1)}{3H} p$$

$$\leq 1 + \frac{2}{3} p$$



Drawback of information: Abilene Topology Case Study



Questions

Thank you for your attention!

Happy to answer questions in the chat forum!

Or by email: Simon Scherrer
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