

PILA: Pervasive Internet-Wide Low-Latency Authentication

Cyrill Krähenbühl Markus Legner Silvan Bitterli
Adrian Perrig
Department of Computer Science
ETH Zurich, Switzerland

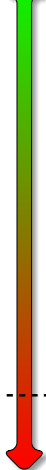
1. July 2021



Motivation

- Trust on first use (TOFU):
 - every on-path entity can attack

Strong Authentication



OWE

TCPCrypt

TOFU

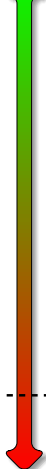
Lowest Level of Security

No Authentication

Motivation

- Trust on first use (TOFU):
 - every on-path entity can attack
 - cannot reliably detect attacks

Strong Authentication



No Authentication

OWE

TCPCrypt

TOFU

Lowest Level of Security

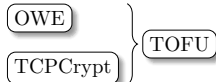
Motivation

- Trust on first use (TOFU):
 - every on-path entity can attack
 - cannot reliably detect attacks
 - cannot pinpoint attacker

Strong Authentication



No Authentication



Motivation

- Trust on first use (TOFU):
 - every on-path entity can attack
 - cannot reliably detect attacks
 - cannot pinpoint attacker
- Strong Authentication:
 - based on PKI (Web PKI or DNSSEC)

Strong Authentication



DANE

Web PKI



OWE

TCPCrypt

TOFU

Lowest Level of Security

No Authentication

Motivation

- Trust on first use (TOFU):
 - every on-path entity can attack
 - cannot reliably detect attacks
 - cannot pinpoint attacker
- Strong Authentication:
 - based on PKI (Web PKI or DNSSEC)
 - name-based authentication

Strong Authentication



DANE

Web PKI



OWE

TCPCrypt

TOFU

Lowest Level of Security

No Authentication

Motivation

- Trust on first use (TOFU):
 - every on-path entity can attack
 - cannot reliably detect attacks
 - cannot pinpoint attacker
- Strong Authentication:
 - based on PKI (Web PKI or DNSSEC)
 - name-based authentication
 - requires configuration

Strong Authentication



DANE

Web PKI



OWE

TCPCrypt

TOFU

Lowest Level of Security

No Authentication

Motivation

- Trust on first use (TOFU):
 - every on-path entity can attack
 - cannot reliably detect attacks
 - cannot pinpoint attacker
- Strong Authentication:
 - based on PKI (Web PKI or DNSSEC)
 - name-based authentication
 - requires configuration
- Can we fill the gap between TOFU and strong authentication?

Strong Authentication



DANE

Web PKI



OWE

TCPCrypt

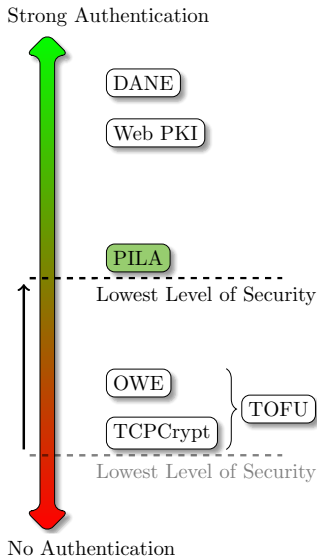
TOFU

Lowest Level of Security

No Authentication

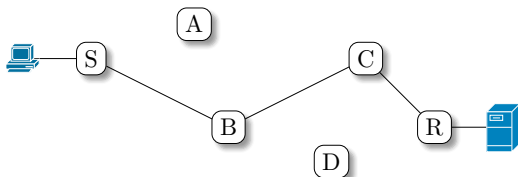
Motivation

- Trust on first use (TOFU):
 - every on-path entity can attack
 - cannot reliably detect attacks
 - cannot pinpoint attacker
- Strong Authentication:
 - based on PKI (Web PKI or DNSSEC)
 - name-based authentication
 - requires configuration
- Can we fill the gap between TOFU and strong authentication?
 - PILA **improves** the base layer for encryption on the Internet



Trust Amplification

- No Authentication



Strong Authentication



DANE

Web PKI

PILA

OWE

TCPCrypt

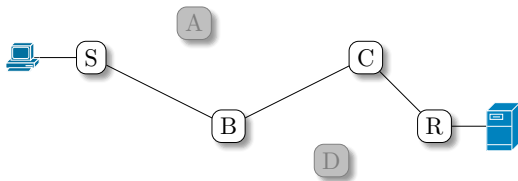
TOFU



No Authentication

Trust Amplification

- No Authentication
- Trust on first use



Strong Authentication



DANE

Web PKI

PILA

OWE

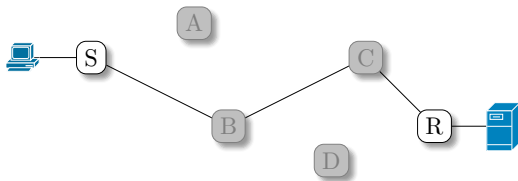
TCPCrypt

TOFU

No Authentication

Trust Amplification

- No Authentication
- Trust on first use
- Trust Amplification
 - Crude Authentication



Strong Authentication

DANE

Web PKI

PILA

OWE

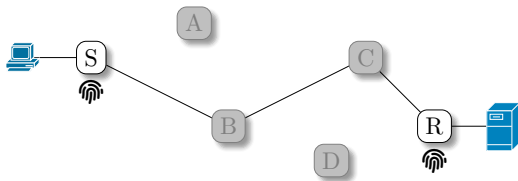
TCPCrypt

TOFU

No Authentication

Trust Amplification

- No Authentication
- Trust on first use
- Trust Amplification
 - Crude Authentication
 - Accountability



Strong Authentication

DANE

Web PKI

PILA

OWE

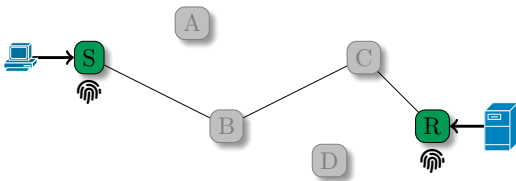
TCPCrypt

TOFU

No Authentication

Trust Amplification

- No Authentication
- Trust on first use
- Trust Amplification
 - Crude Authentication
 - Accountability
 - Leverage



Strong Authentication

DANE

Web PKI

PILA ←

OWE

TCPCrypt

TOFU

No Authentication

Goals

Authentication should ...

- be widely applicable

Goals

Authentication should ...

- be widely applicable
- be low-latency

Goals

Authentication should ...

- be widely applicable
- be low-latency
- require no user interaction

Goals

We propose *PILA*:

Pervasive Internet-Wide Low-Latency Authentication

Authentication should ...

- be widely applicable
- be low-latency
- require no user interaction

Goals

We propose *PILA*:

Pervasive Internet-Wide Low-Latency Authentication

PILA ...

- uses IP-address-based authentication

Authentication should ...

- be widely applicable
- be low-latency
- require no user interaction

Goals

We propose *PILA*:

Pervasive Internet-Wide Low-Latency Authentication

PILA ...

- uses IP-address-based authentication
- has a minimal latency overhead

Authentication should ...

- be widely applicable
- be low-latency
- require no user interaction

Goals

We propose *PILA*:

Pervasive Internet-Wide Low-Latency Authentication

PILA ...

- uses IP-address-based authentication
- has a minimal latency overhead
- automatically generates and fetches certificates

Authentication should ...

- be widely applicable
- be low-latency
- require no user interaction

Goals

We propose *PILA*:

Pervasive Internet-Wide Low-Latency Authentication

PILA ...

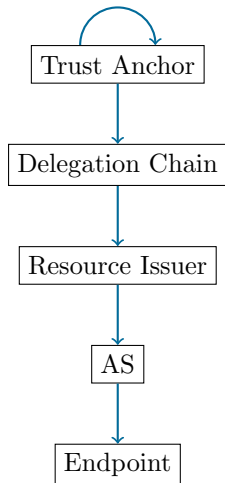
- uses IP-address-based authentication
- has a minimal latency overhead
- automatically generates and fetches certificates
- increases security of TOFU key establishment (only used if strong authentication protocols are not available)

Authentication should ...

- be widely applicable
- be low-latency
- require no user interaction

RPKI as Trust Root

- IANA/RIRs as trust anchor



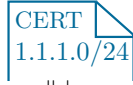
IU



IU



IU

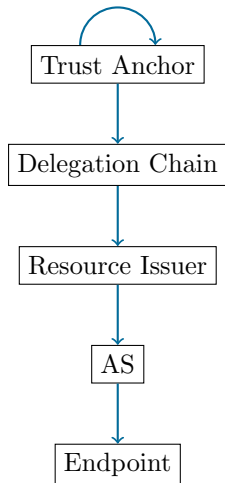


IU



RPKI as Trust Root

- IANA/RIRs as trust anchor
- AS issues short-lived certificates for an IP address to endpoints



CERT
0.0.0.0/0

IU

CERT
1.0.0.0/8

IU

CERT
1.1.0.0/16

IU

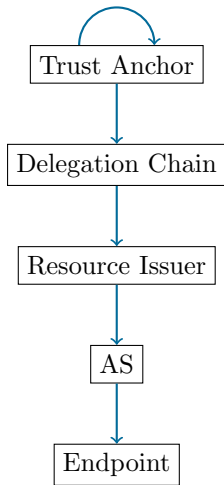
CERT
1.1.1.0/24

IU

CERT
1.1.1.1/32

RPKI as Trust Root

- IANA/RIRs as trust anchor
- AS issues short-lived certificates for an IP address to endpoints
- AS misbehavior (i.e., equivocation) is detectable and cryptographically provable



CERT
0.0.0.0/0

IU

CERT
1.0.0.0/8

IU

CERT
1.1.0.0/16

IU

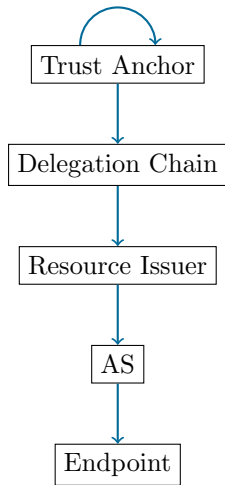
CERT
1.1.1.0/24

IU

CERT
1.1.1.1/32

RPKI as Trust Root

- IANA/RIRs as trust anchor
- AS issues short-lived certificates for an IP address to endpoints
- AS misbehavior (i.e., equivocation) is detectable and cryptographically provable
- ASes are curious but cautious



CERT
0.0.0.0/0

IU

CERT
1.0.0.0/8

IU

CERT
1.1.0.0/16

IU

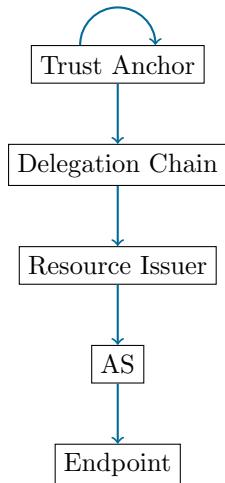
CERT
1.1.1.0/24

IU

CERT
1.1.1.1/32

RPKI as Trust Root

- IANA/RIRs as trust anchor
- AS issues short-lived certificates for an IP address to endpoints
- AS misbehavior (i.e., equivocation) is detectable and cryptographically provable
- ASes are curious but cautious
- Flexible PKI choice (e.g., control-plane PKI in SCION)



CERT
0.0.0.0/0

IU

CERT
1.0.0.0/8

IU

CERT
1.1.0.0/16

IU

CERT
1.1.1.0/24

IU

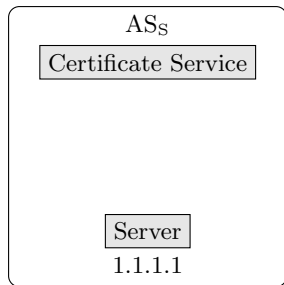
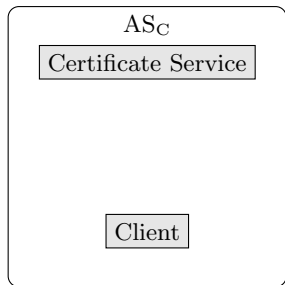
CERT
1.1.1.1/32

Use Cases

- Remote Login (SSH)
- Secure Session-Establishment (TLS)
- Query-Response (DNS)

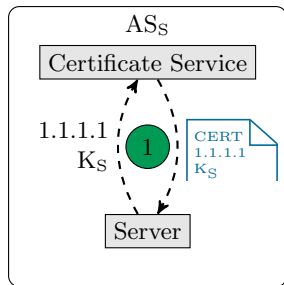
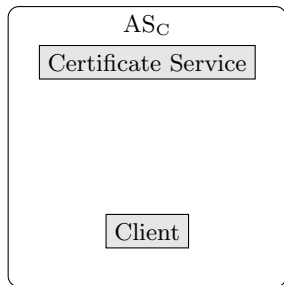
SSH PILA

Server at **1.1.1.1** wants to authenticate itself to the client



SSH PILA

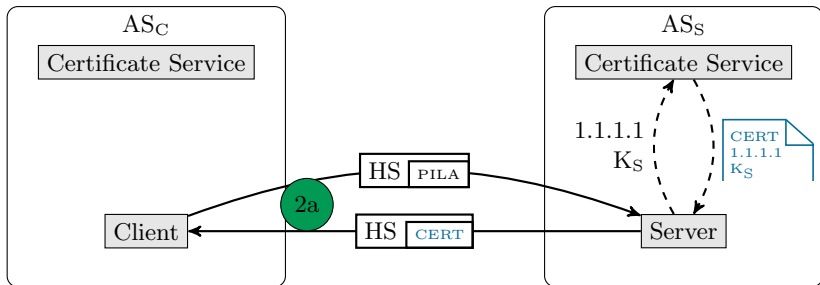
Server periodically fetches short-lived certificate from its local certificate service



SSH PILA

In parallel:

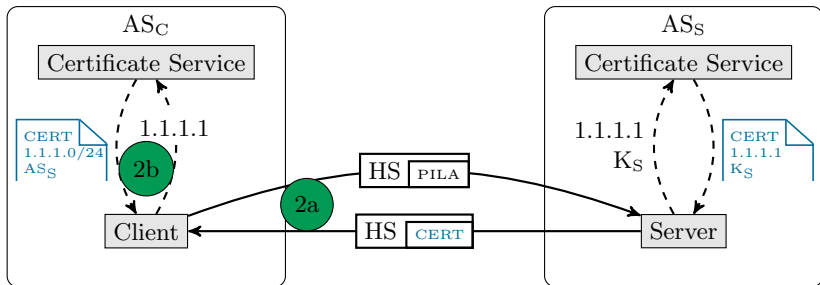
- SSH_{PILA} Handshake (reply contains the certificate)



SSH PILA

In parallel:

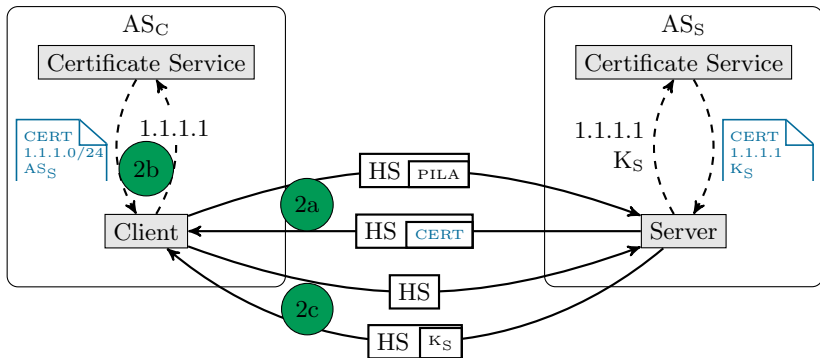
- SSH_{PILA} Handshake (reply contains the certificate)
- Client fetches AS certificate for 1.1.1.1



SSH PILA

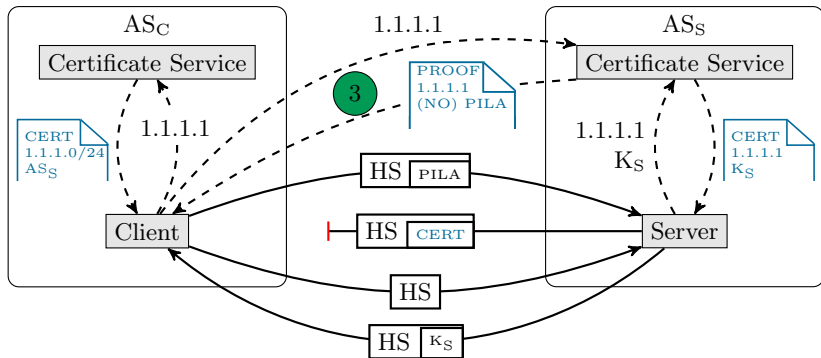
In parallel:

- SSH_{PILA} Handshake (reply contains the certificate)
- Client fetches AS certificate for 1.1.1.1
- Regular SSH Handshake (reply contains the public key)



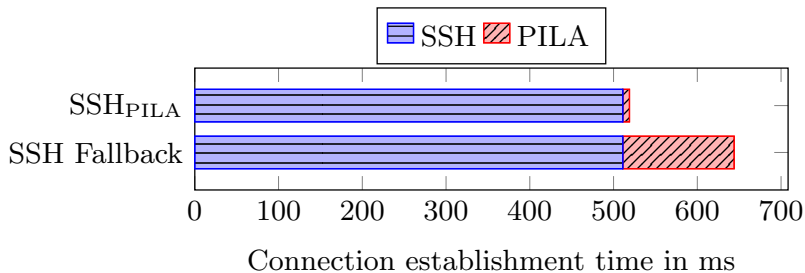
SSH PILA

If the SSH_{PILA} handshake fails, the client requests a proof that the server does not support PILA



SSH PILA

Latency Overhead



SSH PILA

Processing Delay

Average processing times of SSH_{PILA} operations in ms at the client, server, and certificate service:

	Client	Server	Certificate Service
Handshake Overhead	0.8	0.1	-
GetEPCert	-	1.0	17.0
GetASCert	4.3	-	8.3
GetProof	0.6	-	5.1

Conclusion

- Increase security through trust amplification
- PILA offers a new minimum level for fully automatic low latency key establishment
- Implementation and evaluation of PILA in combination with SSH, TLS, and DNS

Thank you!

Cyrill Krähenbühl
Network Security Group
Department of Computer Science
ETH Zürich

cyrill.kraehenbuehl@inf.ethz.ch